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(54) System and method for managing a plurality of snapshots of a file system

(57) A system and method for managing a plurality of snapshots as provided. A set of metadata describing a file system is contained within the file system so that a snapshot of the file system includes the associated metadata. Backup client file systems are restored to a backup server using conventional dump and restore

techniques. The backup server then utilizes a user-defined snapshot management schedule to manage the set of backups associated with the backup server. Such management of snapshots can include deletion of snapshots based upon a variety of parameters including the timestamp.

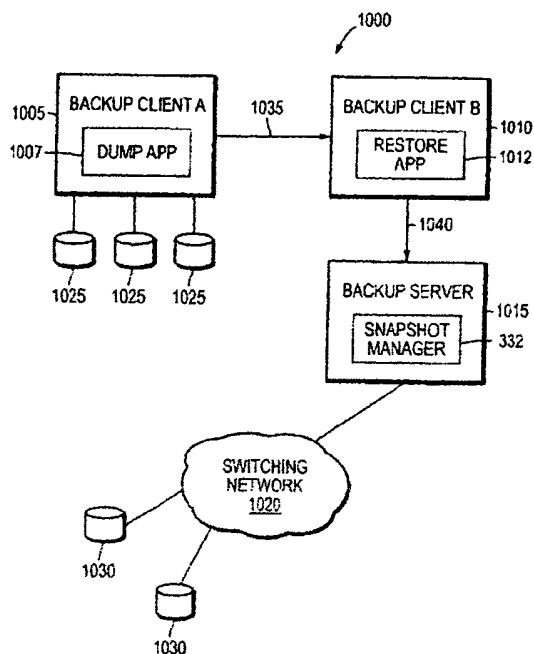


FIG. 10

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Description

FIELD OF THE INVENTION

[0001] The present invention relates to data protection and restoral and, in particular to managing multiple backups of data.

BACKGROUND OF THE INVENTION

[0002] A file server is a computer that provides file service relating to the organization of information on storage devices, such as disks. The file server or filer includes a storage operating system that implements a file system to logically organize the information as a hierarchical structure of directories and files on the disks. Each "on-disk" file may be implemented as a set of disk blocks configured to store information, such as text, whereas the directory may be implemented as a specially-formatted file in which information about other files and directories are stored. A filer may be configured to operate according to a client/server model of information delivery to thereby allow many clients to access files stored on a server, e.g., the filer. In this model, the client may comprise an application, such as a file system protocol, executing on a computer that "connects" to the filer over a computer network, such as a point-to-point link, shared local area network (LAN), wide area network (WAN), or virtual private network (VPN) implemented over a public network such as the Internet. Each client may request the services of the filer by issuing file system protocol messages (in the form of packets) to the filer over the network.

[0003] A common type of file system is a "write in-place" file system, an example of which is the conventional Berkeley fast file system. In a write in-place file system, the locations of the data structures, such as inodes and data blocks, on disk are typically fixed. An inode is a data structure used to store information, such as metadata, about a file, whereas the data blocks are structures used to store the actual data for the file. The information contained in an inode may include, e.g., ownership of the file, access permission for the file, size of the file, file type and references to locations on disk of the data blocks for the file. The references to the locations of the file data are provided by pointers, which may further reference indirect blocks that, in turn, reference the data blocks, depending upon the quantity of data in the file. Changes to the inodes and data blocks are made "in-place" in accordance with the write in-place file system. If an update to a file extends the quantity of data for the file, an additional data block is allocated and the appropriate inode is updated to reference that data block.

[0004] Another type of file system is a write-anywhere file system that does not overwrite data on disks. If a data block on disk is retrieved (read) from disk into memory and "dirtyed" with new data, the data block is stored

(written) to a new location on disk to thereby optimize write performance. A write-anywhere file system may initially assume an optimal layout such that the data is substantially contiguously arranged on disks. The optimal disk layout results in efficient access operations, particularly for sequential read operations, directed to the disks. A particular example of a write-anywhere file system that is configured to operate on a filer is the Write Anywhere File Layout (WAFL™) file system available from Network Appliance, Inc. of Sunnyvale, California. The WAFL file system is implemented within a microkernel as part of the overall protocol stack of the filer and associated disk storage. This microkernel is supplied as part of Network Appliance's Data ONTAP™ storage operating system, residing on the filer, that processes file-service requests from network-attached clients.

[0005] As used herein, the term "storage operating system" generally refers to the computer-executable code operable on a storage system that implements file system semantics and manages data access. In this sense, Data ONTAP software is an example of such a storage operating system implemented as a microkernel. The storage operating system can also be implemented as an application program operating over a general-purpose operating system, such as UNIX® or Windows NT®, or as a general-purpose operating system with configurable functionality, which is configured for storage applications as described herein.

[0006] Disk storage is typically implemented as one or more storage "volumes" that comprise physical storage disks, defining an overall logical arrangement of storage space. Currently available filer implementations can serve a large number of discrete volumes (150 or more, for example). Each volume is associated with its own file system and, for purposes hereof, volume and file system shall generally be used synonymously. The disks within a volume are typically organized as one or more groups of Redundant Array of Independent (or Inexpensive) Disks (RAID). RAID implementations enhance the reliability/integrity of data storage through the redundant writing of data "stripes" across a given number of physical disks in the RAID group, and the appropriate caching of parity information with respect to the striped data. In the example of a WAFL file system, a RAID 4 implementation is advantageously employed. This implementation specifically entails the striping of data across a group of disks, and separate parity caching within a selected disk of the RAID group. As described herein, a volume typically comprises at least one data disk and one associated parity disk (or possibly data/parity) partitions in a single disk) arranged according to a RAID 4, or equivalent high-reliability, implementation.

[0007] In a file server environment, data protection is typically implemented by generating a backup of selected volumes and/or files systems. These backups are generally stored on a tape drive. In certain known file

server configurations, a full backup of the entire file system or volumes is initially created. This full backup stores all of the data contained in the selected volume or file system. At set intervals thereafter, incremental backups are generated. These incremental backups record the changes or *deltas*, between the full backup or last incremental backup and the current state of the data. These backups, both full and incremental, are typically written to a tape drive. A noted disadvantage of writing backups to tape devices is the relatively at which they commit backup data to storage. Overall server performance may be substantially degraded during the backup operation due to the large processing overhead involved with a tape backup operation. This processing overhead derives from copying operations involving the large amount of data that is being moved from the disks comprising the file system or volume to the backup tape device.

[0008] When restoring a file system from a tape backup, many incremental backups are utilized to fully restore the file system. Each of the deltas, or incremental backups, must be individually restored, in the proper order, to generate the active file system. Thus, to fully restore a file system from a set of tape backups, the full backup must first be restored. Then each of the incremental backups, are restored in the proper order to the file system.

[0009] Given the slow speed and other above-noted disadvantages to a tape backup system it seems clear that many administrators would prefer a tapeless backup alternate. One commercially available tapeless backup is the CommVault® Galaxy™ Storage Management Software produced by CommVault Systems of Oceanport, New Jersey. The CommVault system utilizes magneto-optic or compact disc drives in a jukebox setting. This known implementation permits random access of the data so that single file restoration is possible. However, a notable disadvantage to such jukebox systems is that they require a large number of optical disks to contain the entire body of the full and incremental backups. These optical disks need to be stored, handled and changed and are, thus subject to loss or damage.

SUMMARY OF THE INVENTION

[0010] The disadvantages of the prior art are overcome by providing a system and method for managing a plurality of snapshots. To facilitate management, a metadata file is created in the active file system to be snapshotted that stores data concerning the state of the file system. This state information can then be utilized by various application programs to determine how a particular snapshot should be managed. Including this metadata within a snapshot of a file system makes the snapshot become self-describing by including various state information of the snapshotted file system. The resulting system enables a fast and low-overhead tapeless backup at a remote destination backup server.

[0011] The destination backup server, utilizing a file system capable of producing snapshots, is adapted so that backup clients of the backup server can perform conventional backup and recovery or replication operations to replicate backup client file systems on the backup server. At predetermined times, the backup server's file system is snapshotted. A process within the storage operating system of the backup server then manages these collected snapshots according to a user-defined schedule. This management of the plurality of snapshots can include such activities as deletion of stored snapshots based upon a variety of parameters including, e.g., the timestamp stored in the metadata of a snapshot.

[0012] A backup client can selectively restore a path to a desired file without fully restoring an entire file system. This restoring on demand is accomplished by restoring the root directory of a snapshotted file system and then only restoring directories in the path to a particular file. The inodes of the partially restored directories are marked to alert the file system of a backup client to access the backup server for additional files and/or directories.

BRIEF DESCRIPTION OF THE DRAWINGS

[0013] The above and further advantages of the invention may be better understood by referring to the following description in conjunction with the accompanying drawings in which like reference numerals indicate identical or functionally similar elements:

Fig. 1 is a schematic block diagram of an exemplary network environment having file servers and tape drives;

Fig. 2 is a schematic block diagram of an exemplary file server in accordance with an embodiment of this invention;

Fig. 3 is schematic block diagram of an exemplary storage file system for use in accordance with an embodiment of this invention;

Fig. 4 is schematic block diagram of an exemplary file system inode structure;

Fig. 5 is a schematic block diagram of the exemplary file system inode structure of Fig. 4 including a snapshot inode;

Fig. 6 is a schematic block diagram of an exemplary file system inode structure of Fig. 4 after data block has been rewritten;

Fig. 7 is a time line of a stable and transitioning file system;

Fig. 8 is a schematic block diagram of snapshot metadata in accordance with an embodiment of this invention;

Fig. 9 is a flow chart detailing the procedure performed in utilizing stored metadata to select a proper snapshot;

Fig. 10 is a schematic block diagram of an exem-

plary networking environment including file servers in accordance with an embodiment of this invention; Fig. 11 is a flow chart detailing the procedure utilized by file servers in the illustrative environment of Fig. 10 in accordance with an embodiment of this invention;

Fig. 12 is a schematic block diagram of an exemplary directory structure;

Fig. 13 is a schematic block diagram of a partially restored directory structure in accordance with an embodiment of the invention; and

Fig. 14 is a schematic block diagram of an exemplary inode.

DETAILED DESCRIPTION OF AN ILLUSTRATIVE EMBODIMENT

A. Network Environment

[0014] Fig. 1 is a schematic block diagram of an exemplary network environment 100 in which the principles of one embodiment of the present invention are implemented. A network 100 is based around a network cloud 102. This network cloud can be a local or network (LAN), a wide area network (WAN), virtual private network (VPN) utilizing communication links over the internet, for example, or a combination of LAN, WAN and VPN implementations can be established. For the purposes of this description, the term network cloud should taken broadly to include any acceptable network architecture. The network cloud 102 interconnects various clients 104. Also attached to the network cloud is a file server 200. This file server, described further below is configured to control storage of, and access to, data and a set 108 of interconnected storage volumes 106. Connected to the file server is a tape drive 110. In known data backup examples, the file server 200 would back up data stored on disks 106 to the tape drive 110. Each of the devices attached to the network cloud include an appropriate conventional network interface arrangement (not shown) for communicating over the network cloud using desired communication protocols such as the well-known Transport Control Protocol/Internet Protocol (TCP/IP), User Datagram Protocol (UDP), Hyper Text Transport Protocol (HTTP), Simple Network Management Protocol (SNMP), or Virtual Interface Connections (VI).

B. File Servers

[0015] Fig. 2 is a more-detailed schematic block diagram of an exemplary file server 200. By way of background, a file server or *filer*, is a computer that provides file service relating to the organization of information on storage devices, such as disks. However, it will be understood by those skilled in the art that the inventive concepts described here may apply to any type of file server, wherever implemented as a special-purpose or general-

purpose computer, including a standalone computer.

[0016] The file server 200 comprises a processor 202, in memory 204, in network adapter 206, a nonvolatile random access memory (NVRAM) 208 in the storage adapter 210 interconnected by system bus 212. Contained within the memory 204 is a storage operating system 300 that implements a file system to logically organize the information as a hierarchical structure of directories and files on the disks. In the illustrative embodiment, the memory 204 comprises storage locations that are addressable by the processor and adapters for storing software program code. The operating system 300, portions of which are typically resident in memory and executed by the processing elements, functionally organizes the filer by *inter alia*, invoking storage operations in support of a file service implemented by the file server.

[0017] The network adapter 206 comprises a mechanical, electrical and signaling circuitry needed to connect the file server 200 to client 104 over network cloud 102. The client 104 maybe a general-purpose computer configured to execute applications, such as data base applications. Moreover, the client 104 may interact with the filer server 200 in accordance with the client/server model of information delivery. That is, the client may request the services of the file server, and the file server may return the results of the services requested by the client, by exchanging packets defined by an appropriate networking protocol.

[0018] The storage adapter 210 incorporates with the storage operating system 300 executing on the file server to access information requested by the client. Information maybe stored on the disks 106 of a disk 108 (Fig. 1) that is attached via the storage adapter 210 to the file server. The storage adapter 210 includes input/output (I/O) interface circuitry that couples to the disks over in I/O interconnect arrangement, such as a conventional high-performance Fibre Channel serial link topology. The information is retrieved by the storage adapter and, if necessary, processed by the processor 202 (or the adapter 210 itself) prior to be forwarded over the system bus 212 to the network adapter 206, where information is formatted into appropriate packets and returned to the client 104.

[0019] In one exemplary file server implementation, the file server can include a nonvolatile random access memory (NVRAM) 208 that provides fault-tolerant backup of data, enabling the integrity of filer server transactions to survive a service interruption based upon a power failure, or other fault.

C. Storage Operating System

[0020] To facilitate the generalized access to the disks 106 on the array 108, the storage operating system 300 implements write-anywhere file system that logically organizes the information as a hierarchical structure of directories and files on the disks. Each "on-disk" file may

be implemented as a set of disks blocks configured to distort information, such as data, where as the directory may be implemented as a specially formatted file which other files and directories are stored. As noted above, in the illustrative embodiment described herein, the operating system is the NetApp® Data ONTAP™ operating system available from Network Appliance, Inc., that implements the write-anywhere file layout (WAFL™) file system. It is expressly contemplated that any appropriate file system can be used, and as such, where the term WAFL or file system is employed, it should be taken broadly to refer to any file system that is otherwise adaptable to the teachings of this invention.

[0021] The storage operating system comprises a series of software layers, including a media access layer 302 of network drivers (e.g., an Ethernet driver). The storage operating system 300 further includes network protocol layers, such as an Internet Protocol (IP) layer 304 and its supporting transport mechanisms, the Transport Control Protocol (TCP) layer 306 and the User Datagram Protocol (UDP) layer 308.

[0022] A file system protocol layer provides multi-protocol data access and, to that end, includes support for the Network File System (NFS) protocol 312, the Common Internet File System (CIFS) protocol 314 and the Hyper Text Transfer Protocol (HTTP) 316. In addition, the storage operating system 300 includes a disk storage layer 322 that implements a disk storage protocol, such as a RAID protocol, and a disk driver layer 324 that implements a disk access protocol such as, e.g., a Small Computer System Interface (SCSI) protocol.

[0023] Bridging the disk software layers with the network and file system protocol layers is a file system layer 326 of the storage operating system 300. Generally the file system layer 326 implements a file system having an on-disk format representation that is block-based using, e.g., 4-kilobyte (KB) data blocks and using inodes to describe the files. In response to transaction requests, the file system generates operations to load (retrieve) the requested data from disks 106 if it is not resident "in-core", i.e., in the filer's memory 204. If the information is not in memory, the file system layer 326 indexes into the inode file using the inode number to access an appropriate entry and retrieve a logical volume block number. The file system layer 326 then passes the logical volume block number to the disk storage (RAID) layer 322, which maps that logical number to a disk block number and sends the latter to an appropriate driver (for example, an encapsulation of SCSI implemented on a fibre channel disk interconnection) of the disk driver layer 324. The disk driver accesses the disk block number from disks 106 and loads the requested data in memory 204 for processing by the filer 200. Upon completion of the request, the filer (and storage operating system) returns a reply, e.g., a conventional acknowledgement packet defined by the CIFS specification, to the client 104 over the network cloud 102.

[0024] It should be noted that the storage access re-

quest data path 330 through storage operating system layers described above needed to perform data storage access for the client requests received the file server may alternately be implemented in hardware, software or a combination of hardware and software. That is, in an alternative embodiment of this invention, the storage access request data path 330 may be implemented as logic circuitry embodied within a field programmable gate array (FPGA) or in an application specific integrated circuit (ASIC). This type of hardware implementation increases the performance of the file service provided by the file server 200 in response to a file system request issued by a client.

Included within the file system layer is a set of snapshot processes 328, which implement the inherent snapshot capabilities of the file system. The native Snapshot™ capabilities of the WAFL file system are further described in *TR3002 File System Design for an NFS File Server Appliance* by David Hitz et al., published by Network Appliance, Inc., and in U.S. Patent No. 5,819,292 METHOD FOR MAINTAINING CONSISTENT STATES OF A FILE SYSTEM AND FOR CREATING USER-ACCESSIBLE READ-ONLY COPIES OF A FILE SYSTEM by David Hitz et al. "Snapshot" is a trademark of Network Appliance, Inc. It is used for purposes of this patent to designate a persistent consistency point (CP) image. A persistent consistency point image (PCPI) is a point-in-time representation of the storage system, and more particularly, of the active file system, stored on a storage device (e.g., on disk) or in other persistent memory and having a name or other unique identifier that distinguishes it from other PCPIs taken at other points in time. A PCPI can also include other information (metadata) about the active file system at the particular point in time for which the image is taken. The terms "PCPI" and "snapshot" shall be used interchangeably through out this patent without derogation of Network Appliance's trademark rights.

[0025] By way of background, a snapshot is a restorable version of a file system created at a predetermined point in time. Snapshots are generally created on some regular schedule. The snapshot is stored on-disk along with the active file system, and is called into the buffer cache of the filer memory as requested by the storage operating system. An exemplary file system inode structure 400 is shown in Fig. 4. The root inode 405 contains information describing the inode file associated with a given file system. In this exemplary file system inode structure root inode 405 contains a pointer to the inode file indirect block 410. The inode file indirect block 410 contains a set of pointers to inode file and data blocks 415. The mode file data block 415 includes pointers to file and data blocks to 420A, 420B and 420C. Each of the file data blocks 420(A-C) is capable of storing, in the illustrative embodiment, 4 kilobytes (KB) of data.

[0026] When the file system generates a snapshot of a given file system, a snapshot inode is generated as shown in Fig. 5. The snapshot inode 505 is, in essence,

a duplicate copy of the root inode 405 of the file system 400. Thus, the exemplary file system structure 500 includes the same inode file indirect block 410, inode file data block(s) 415 and file data blocks 420A-C as in Fig. 4. When a user modifies a file data block, the file system layer writes the new data block to disk and changes the active file system to point to the newly created block.

[0027] Fig. 6 shows an exemplary inode file system structure 600 after a file data block has been modified. In this illustrative example, file data block 420C was modified to file data block 420C'. When file data block 420C is modified to file data block 420C', the contents of the modified file data block are written to a new location on disk as a function for the exemplary WAFL file system. Because of this new location, the inode file data block 615 pointing to the revised file data block 420C must be modified to reflect the new location of the file data block 420C. Similarly, the inode file indirect block 610 must be rewritten to point to the newly revised inode file and data block. Thus, after a file data block has been modified the snapshot inode 505 contains a point to the original inode file system indirect block 410 which in turn contains a link to the inode file data block 415. This inode file data block 415 contains pointers to the original file data blocks 420A, 420B and 420C. However, the newly written inode file data block 615 includes pointers to unmodified file data blocks 420A and 420B. The inode file data block 615 also contains a pointer to the modified file data block 420C' representing the new arrangement of the active file system. A new file system root inode 605 is established representing the new structure 600. Note that metadata (not shown) stored in any snapshotted blocks (e.g., 505, 410, and 420C) protects these blocks from being recycled or overwritten until they are released from all snapshots. Thus, while the active file system root inode 605 points to new blocks 610, 615 and 420C', the old blocks 505, 410, 415 and 420C are retained until the snapshot is fully released.

[0028] After a snapshot has been created and file data blocks modified, the file system layer can reconstruct or "restore" the file system inode structure as it existed at the time of the snapshot by accessing the snapshot inode. By following the pointers contained in the snapshot inode 505 through the inode file indirect block 410 and inode file data block 415 to the unmodified file data blocks 420A-C, the file system layer can reconstruct the file system as it existed at the time of creation of the snapshot.

[0029] In alternate embodiments of a storage operating system, the snapshot management layer 332 permits users to define a set schedule of operations to be performed on a plurality of snapshots. The user defines the set schedule by inputting via a command line interface (CLI) or a graphical user interface (GUI) specific operations and parameters to be performed on the snapshots created by the snapshot process 328. Typically, backup servers, described further below, would execute a storage operating system including a snap-

shot management layer. However, it is expressly contemplated that alternate implementations of file servers or computers could execute a storage operating system including the functionality of a snapshot management layer.

D. Snapshot Contents

[0030] As snapshots of a file system can be generated at any time, problems can occur if a snapshot is taken of a file system while it is being updated. An exemplary timeline 700 of the states that a file system undergoes is shown in Fig. 7. The file system is initially stable during time 705. By "stable" it is meant that no updates are being physically written to the disks or inode structures comprising the file system. The file system then enters a transitioning phase 715 before reaching another instance of a stable file system 710. During the transitioning phase 715, data is written to the file system or the file system is otherwise modified, e.g., files are deleted. Thus, while the file system is transmitting, particular inodes and data blocks are modified. As noted above, snapshots of a file system can be generated at any time. Thus, for example, a snapshot, could be generated at time 720 (stable period) or 725 (transitioning period). A snapshot generated at time 720, would capture a stable file system with no updates occurring.

[0031] As numerous snapshots are generated at arbitrary periods, some will capture stable file systems and others will capture file systems in transition. For restoration purposes, it is desired to know the status of the file system contained within a particular snapshot. A database could be created and maintained to track each snapshot of a file system and its associated status (i.e., stable or transitioning). However, such a database would require additional computational overhead.

[0032] To enable a user or administrator to determine if a particular snapshot is of a stable or transitioning file system, the state of the file system is stored in the active file system. By "active file system" it is meant generally the file system associated with a particular computer to which current input/output operations are being directed. This metadata, or data associated with the state of the file system, can be stored in a variety of places in the active file system. In one embodiment, a file is created in the active file system to store the metadata. The contents of the metadata stored in the active file system is shown in Fig. 8. Stored within the active file system 800, is a set of metadata 805. This metadata 805 includes various entries, including entries for the state 810, the source filer 815, a time stamp 820 and, in alternate embodiments, additional entries 825. The state entry 810 can be implemented as a binary flag. If, for example, the value of the state flag is "1", the file system is stable. Alternately, if the value of the state flag is "0", the file system is in a transitioning state. Thus, the file system can modify the state entry to show that the file system is transitioning before the file system begins to

modify data within the file system. Thus, any snapshots taken during the course of the data update, will include the metadata showing that the file system is in a transitioning state. Once the data update is complete, a file system modifies the state entry 810 to show that the file system is stable. In one embodiment, this metadata is stored in a qtree metafile.

[0033] The source filer entry 815 is used to identify the file server originating a particular file system. Such a flag can be useful when, for example, a plurality of snapshots from differing file servers are stored on a single volume.

[0034] By utilizing the time stamp entry 820 in the metadata 805, a file server can determine if a particular snapshot contains the desired version of the file system. If, for example, a user desired to restore a file or files from a particular time frame, the user or the administrator would need to know the particular time a snapshot was created. Fig. 9 is a flow chart detailing procedure performed by a user in determining if a particular snapshot is a proper snapshot for a desired operation (e.g., restoring a file from a certain point-in-time). In step 905, the user selects the first snapshot from a list of available snapshots. This list of snapshots is maintained by the file system. Next, in step 910, the process would look at the metadata file in the selected snapshot. By sorting or analyzing the data contained within the metadata a determination can be made if this snapshot is the proper snapshot (step 915). If the snapshot is the proper snapshot, the desired operation is performed on/with the selected snapshot (step 920). If the snapshot is not the proper snapshot, then the process selects the next snapshot from a list of available snapshots (step 925). After selecting the next snapshot from the list, the process branches to step 910 to query the metadata contained within the selected snapshot.

[0035] Thus, by storing metadata relating to a snapshot in the active file system, a snapshot of that file system naturally and invariably includes the metadata. In effect, the snapshot becomes self-describing based upon the incorporated metadata. This facilitates management of a plurality of snapshots. Additionally, other process and programs can utilize the metadata stored in the snapshot to make management decisions with regard to the particular snapshot.

E. Snapshot Management

[0036] In accordance with an embodiment of the present invention, to manage a set of backups or snapshots of a series of backup clients, a separate backup server is utilized. Such a backup server can be implemented as, e.g., a filer server executing a storage operating system having a snapshot management layer 332 (Fig. 3). Likewise, any or all of the backup clients can be implemented in accordance with the above description.

[0037] As a backup client of the backup server gen-

erates various backups of the backup client's file systems, these backups are restored to a file system associated with the backup server. Thus, at any given time, the backup server contains a complete file system of an associated backup client. By "backup client" it is meant any other device or computer that utilizes the capabilities of the back up server for storage and management of snapshots. The backed up data of the backup client is not stored as a base line backup along with a series of incremental backups. Instead, the entire file system is replicated on the backup server.

[0038] The snapshot management software 332 can be used to maintain the set of backups generated. An example of such management is deleting certain snapshots over time. It is generally desirable for the backup server to delete snapshots after a certain time period. It is infeasible to store every snapshot generated, as it would require an ever-increasing amount of storage space. The user can define a schedule of storing and maintaining the snapshots stored on the backup server. Thus, the backup client can generate a backup or snapshot at regular intervals. The snapshots are then transferred to the backup server at a predetermined time interval. The backup client is then free to subsequently delete these various snapshots to free disk space on the backup client. However, the backup server manages the set of snapshots in accordance with the user-defined schedule.

[0039] A notable feature of the use of a backup server is that the system administrator is not limited by the tape infrastructure. The backup server can be located at a significantly greater distance (e.g., any distance served by a network) than is possible with a small computer system interface (SCSI) buss commonly attached tape drive. Another advantage of the use of a backup server is that a backup client can use known dump and restore backup software and/or procedures. Such dump and restore backup procedures, in connection with a backup server able to implement snapshots, permits snapshots to be effectively generated of file systems managed by computers which utilize a file system that does not contain inherent snapshot capabilities.

[0040] An example of such a backup network arrangement is shown in Fig. 10. The network environment 1000 includes backup client A 1005, backup client B 1010, and the backup server 1015. Connected to the backup server 1015 is a switching network 1020. A plurality or set of disks 1030 is connected to the switching network. The backup server utilizes these disks to store its file system. Connected to backup client A 1005 is a set of locally connected disks 1025. If backup client A 1005 is running a conventional backup software (Dump App) 1007 which permits a dump stream to be generated, similar to a conventional backup operation, over data path 1035 to server B. Data path 1035 can be any form of data path, including a local area network (LAN), a wide area network (WAN), or other acceptable networking architecture. Backup client B 1010 accepts the

dump stream generated from backup client A 1005 and utilizing conventional restore software (Restore App) 1012 performs a restore operation to the backup server over data path 1040 via a restore stream. A restore stream comprises a series of write operations directed at the backup server's file system. The stream can be encapsulated in a variety of network protocols including TCP/IP.

[0041] The data path 1040, similarly to data path 1035 can be of any acceptable networking architecture. The backup client B 1010 generates the restore stream, which recreates a file system on a volume associated with the backup server. The backup server 1015 can then generate a series of snapshots of the data stored on the volume. As the file system contained on the backup servers disks 1030 is a restored version of the file system stored on backup client A's disks 1025, the backup server, by generating snapshots, is effectively generating snapshots of the file system contained on disks 1025.

[0042] An exemplary procedure performed by such a network environment is shown in Fig. 11. In step 1105, the file server whose file system is to be backed up performs a baseline dump (step 1105). This baseline dump can be generated using any acceptable backup and recovery or replication software and/or hardware. Next, in step 1110, a piped baseline restore is generated on the backup server. This piped baseline restore is generated using the dump data stream generated in the previous step. Once the baseline piped restore has been generated, the backup server creates a snapshot of the file system generated (step 1115). The backup server then waits a predetermined period of time (step 1120). Then, in step 1125, the file server or computer being whose file system is being backed up performs an incremental dump piped to an incremental restore to the backup server. Once this incremental restore has been generated on the backup server, the backup server creates a new snapshot in step 1130. The backup server then manages the set of snapshots according to the end user schedule in step 1135. This end user schedule can include a variety of user defined options. For example, a user may desire to store all snapshots generated in the last 24 hours but only one snapshot per day for the prior week.

[0043] The snapshot management software 332 (Fig. 3), using the user-defined management schedule, performs various operations on the set of snapshots generated. Exemplary operations could include, as noted above, deletion of snapshots which are generated prior to a specified date, deletion of snapshots of the file system which are in a transitioning state, or deletion of snapshots according to any other user-defined methodology. For example, the snapshot management software could be configured so that only the last snapshot generated in each day is retained. Thus, in this example, the snapshot manager would determine if a plurality of snapshots shared a timestamp evidencing that they

were generated on the same day. If a plurality of snapshots were generated on the same day, the snapshot management software would delete all those which were not the last snapshot generated on that day.

F. Restore on Demand

[0044] The backup file system stored on the backup server is complete file system and is randomly accessible as the backup file system is stored on disks unlike serial-based devices such as tape devices. To improve restoration performance, individual files and directories can be restored on demand instead of waiting until the complete file system is restored.

[0045] Fig. 12 is a block diagram of an exemplary directory structure 1200. The directory structure 1200 includes a root directory 1205 having two subdirectories *foo* 1210 and *bar* 1215. Under the *foo* directory 1210 is a subdirectory *sub* 1220. The *sub* subdirectory 1220 includes files A and B 1225 and 1230. Similarly the *bar* directory 1215 includes file C 1235.

[0046] If a backup client of a backup server requires a restoration of a file system, for example the illustrative file system of Fig. 12, the various files and directories can be restored on demand to improve restoration performance. For example, assume that the file system of Fig. 12 was corrupted on the backup client. If the backup server contains a snapshot of the file system, a restoration can be generated from the snapshot. By only restoring on demand those files or directories that are needed, time and processing overhead can be saved. For example, assume that the file 1235 is required by the backup client.

[0047] Under traditional snapshot restoration techniques, the entire file system 1200 would be restored. If, for example, the file system 1200 contained two terabytes (TB) of data, the entire 2 TB of data would need to be copied from the backup server snapshot to the active file system of the backup client. Such a data transfer could take a considerable amount of time. In order to alleviate such need for transferring mass quantities of data, only the file system structures needed to access a particular file are restored. Thus, when file 1235 is requested by the backup client, only the *bar* directory 1215 and the file 1235 need to be restored. Such a restoration is shown in Fig. 13. This restored file system structure 1300 includes the root directory 1205, the *bar* directory 1215 and the required file 1235. Thus, the *foo* directory 1210 and its associated subdirectories and files are not restored unless and until they are requested by the backup client. In an alternate embodiment, a full restoration can be performed using background processing while the restore on demand is being utilized.

[0048] Each inode of the restored file system has a flag that is set that alerts the file system layer of the storage operating system of the backup client that it is a partially restored inode. Thus, in response to such a flag, the file system layer can direct file access requests to

the backup server to continue to restore on demand needed files.

[0049] An exemplary inode 1400 is shown in Fig. 3. The exemplary inode 1400 includes a metadata section 1405 and a data section 1435. The information stored in the metadata section 1405 of each inode 1400 describes a file, and such, includes the type (e.g., regular or directory) 14010 of the file, the size 1415 of a file, time stamps (e.g. access and/or modification) 1420 for the file and ownership, i.e., user identifier (UID 1425) of the file. The metadata section 1405 also includes a flag identifying wherever the inode has been partially restored from a snapshot. The partially restored flag 1430 alerts the file system that the contents of the inode may not be completely restored from a snapshot and, to ensure that the proper data is obtained, the file system should access the backup server snapshot. The contents of the data section 1435 of the inode however, may be interpreted differently depending upon the type of file (inode) defined within the type field 1410. For example, the data section 1435 of a directory inode contains metadata controlled by the file system, whereas the data section of a regular inode contains user-defined data. In this latter case, the data section 1435 includes a representation of the data associated with the file.

[0050] Thus, by only restoring those files and directories as they are needed by the storage operating system, substantial processing time is saved. This time savings can be especially evident when the file system contains a substantial amount of data, but only a few smaller files are routinely needed. In such cases, the regularly accessed files can be restored as needed, while the larger and less frequently used files can be restored using a background process.

[0051] Accordingly, by keeping a metadata file within the active file system, data relating to the file system can be stored. By utilizing the file system's inherent snapshot capabilities, this metadata file will be incorporated into any snapshots of the file system. Thus, a file system becomes, in effect, a self-describing snapshot. The metadata can also be utilized by a snapshot management program running on a backup server. The backup server can accept conventional restore data streams from computers utilizing file systems which do not incorporate snapshot capabilities. Once a backup client of the backup server has restored the backup clients file system to the backup server, the backup server can take snapshots of the restored file system. This allows for the generation of snapshots of file systems which do not inherently contain the capability to generate a snapshot. In addition, this system and method advantageously enables a reliable, fast and low-overhead tapeless backup using a remote destination backup file server.

[0052] The foregoing has been a detailed description of illustrative embodiments of the invention. Various modifications and additions can be made without departing from the scope of the invention. It is expressly contemplated that any of the functions, procedures or

processes described herein can be implemented using hardware, firmware or software, consisting of a computer-readable medium including program instructions executing on a computer, or a combination of hardware, firmware and/or software. The computer readable medium can be any carrier medium such as a signal e.g. an electrical, optical, microwave, magnetic, acoustic or electromagnetic signal, or a storage medium such as a floppy disk, hard disk, CD ROM or solid state memory device. Accordingly, this description is meant to be taken only by way of example and not to otherwise limit the scope of the invention.

15 Claims

1. A method of managing a plurality of snapshots of a file system associated with a backup server, the method comprising the steps of:

creating a baseline snapshot of the file system on the backup server;
performing an incremental restore to the backup server;
creating a new snapshot on the backup server; and
managing a plurality of snapshots, including the baseline and new snapshot, according to a retention policy.

2. The method of claim 1 further comprising the step of waiting a predetermined period of time before creating the new snapshot on the backup server.

3. The method of claim 1 or claim 2, wherein the retention policy is user-defined and specifies a retention period.

4. The method of any preceding claim, wherein the step of performing an incremental restore comprises the step of performing an incremental restore from a set of backup clients to the backup server.

5. The method of any preceding claim, wherein the step of creating a baseline snapshot on the backup server further comprises the steps of:

performing a baseline dump of a file system associated with a backup client, the baseline dump being piped to a baseline file system on the backup server; and
creating a snapshot of the baseline file system on the backup server.

6. The method of any preceding claim, wherein the step of managing a plurality of snapshots further comprises the step of deleting one or more snapshots.

7. The method of claim 6, wherein the step of deleting one or more snapshots of the plurality of snapshots occurs according to the retention policy.
8. A method of managing a plurality of snapshots of a logical group of data blocks associated with a secondary system, the method comprising the steps of:
- creating a baseline snapshot of the logical group on the secondary system;
 - performing an incremental restore from at least one primary system to the secondary system;
 - creating a new snapshot of the logical group associated with the secondary system; and
 - managing the plurality of snapshots, including the baseline and new snapshots, according to a retention policy.
9. The method of claim 8, wherein the step of creating the baseline snapshot further comprises the steps of:
- replicating a backup client file system in the file system associated with the backup server; and
 - creating a baseline snapshot of the file system associated with the backup server.
10. The method of claim 8 or claim 9, wherein the step of managing the plurality of snapshots further comprises deleting one or more of the plurality of snapshots in accordance with the retention policy.
11. A backup server for operative interconnection with at least one backup client, the backup server comprising:
- means for creating a replica of a file system associated with the at least one backup client on one or more file systems associated with the backup server;
 - means for creating a snapshot of the one or more file systems associated with the backup server;
 - means for incrementally updating the one or more file systems associated with the backup server in response to an incremental restore operation from the one or more backup clients; and
 - means for managing a plurality of snapshots in accordance with a retention policy.
12. The backup server of claim 11, wherein the retention policy is user-defined and specifies a retention period.
13. A method of generating a self-describing snapshot of a file system, the method comprising the steps of:
- creating a metadata file in the file system;
 - storing file system information in the metadata file; and
 - generating a snapshot of the file system, wherein the metadata file is included in the snapshot of the file system.
14. The method of claim 13, wherein the file system information further comprises state information associated with the file system.
15. The method of claim 14, wherein the state information identifies the file system as being in a stable state.
16. The method of claim 14, wherein the state information identifies the file system as being in a transitioning state.
17. The method of any one of claims 14 to 16, wherein the file system information further comprises a time stamp.
18. A method of accessing files stored in a snapshot of a file system, the method of comprising the steps of:
- restoring a root directory of the file system;
 - restoring a set of directories, the set of directories being associated with a path to the file; and
 - restoring the file from the snapshot, wherein directories not associated with the path to the file are not restored.
19. The method of claim 18, wherein the set of restored directories further comprises a set of inodes, the set of inodes including information identifying the set of inodes as being partially restored.
20. A computer readable medium, carrying program instructions for controlling a computer to carry out the method of any one of claims 1 to 10, or 13 to 19.

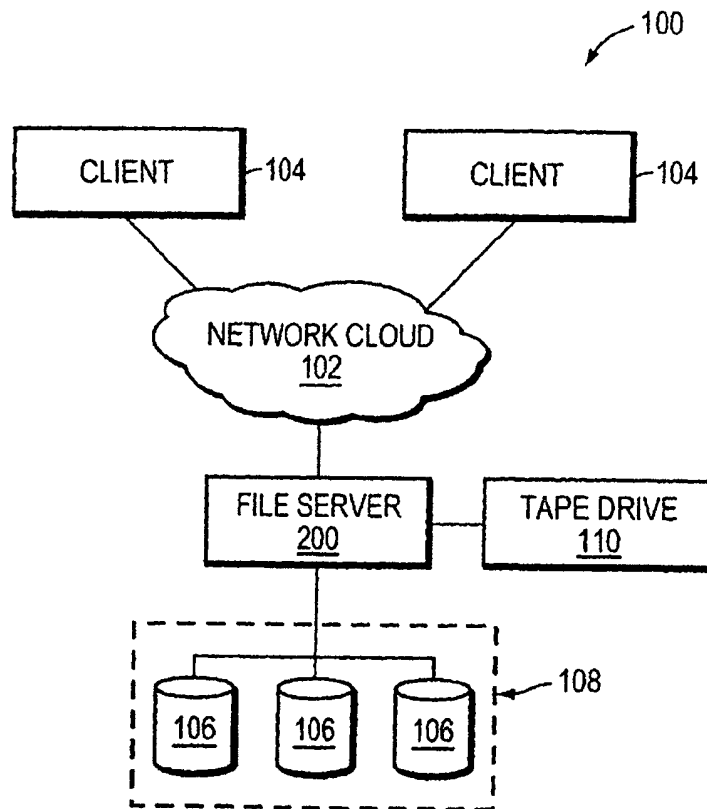


FIG. 1

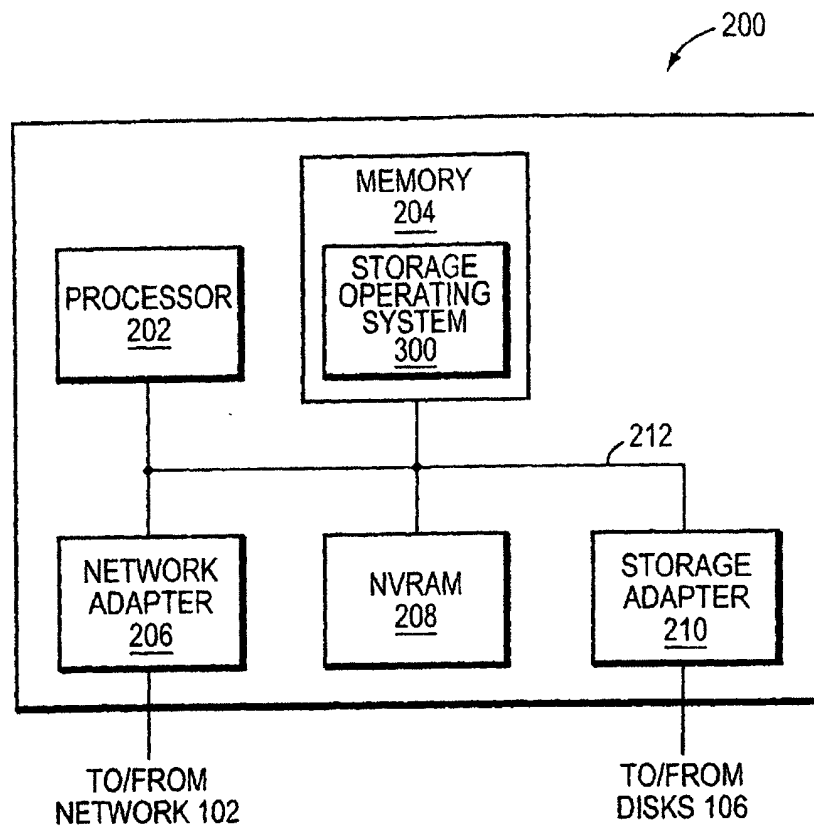


FIG. 2

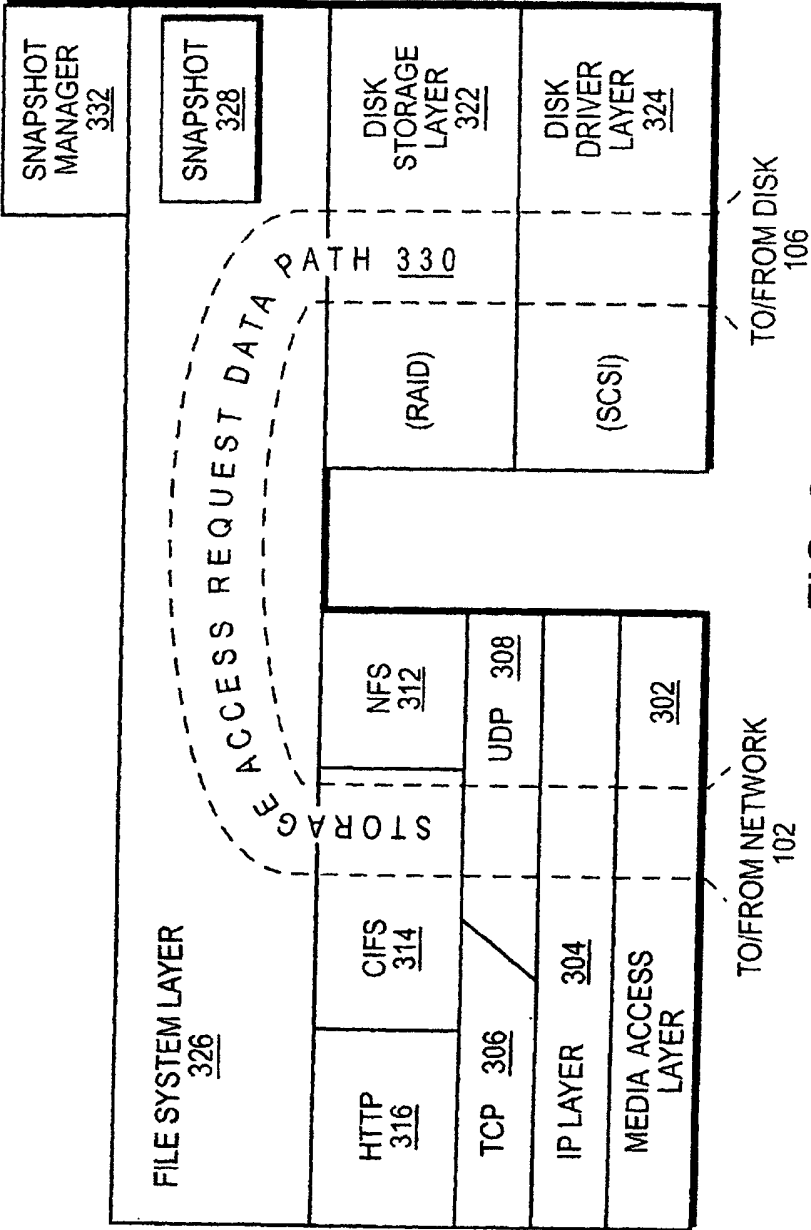


FIG. 3

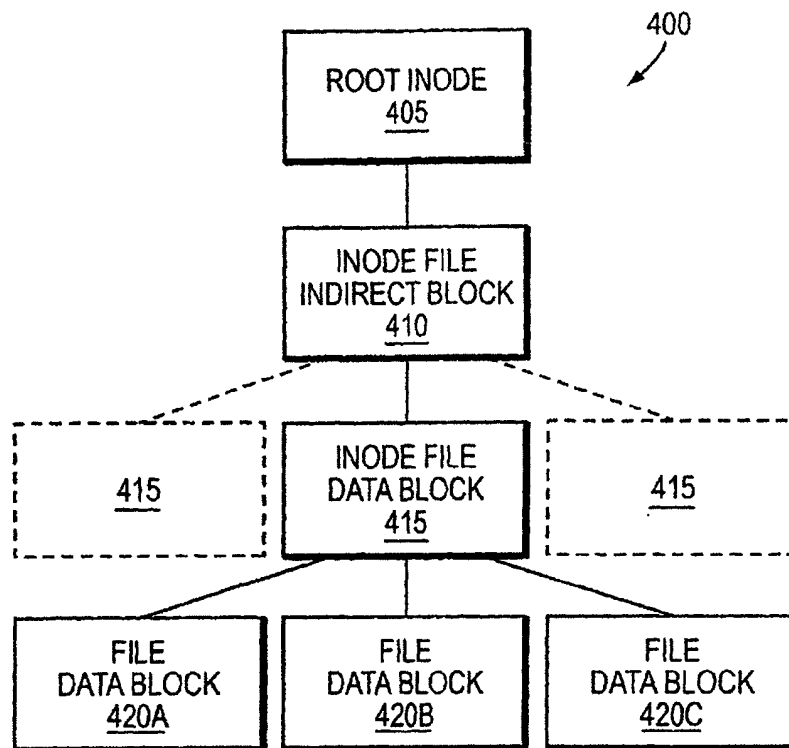


FIG. 4

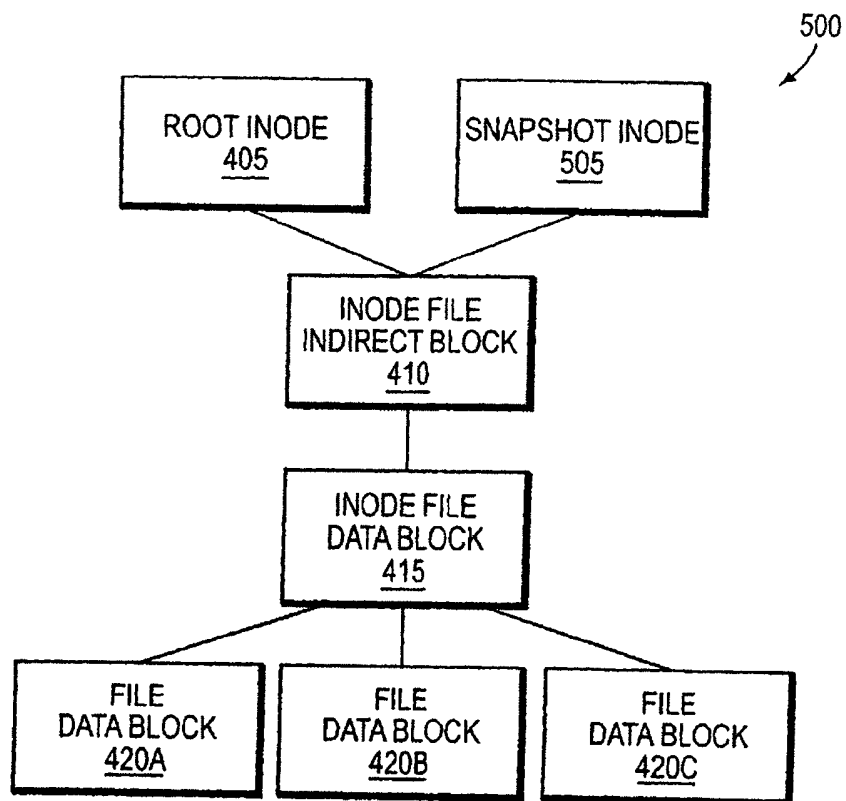


FIG. 5

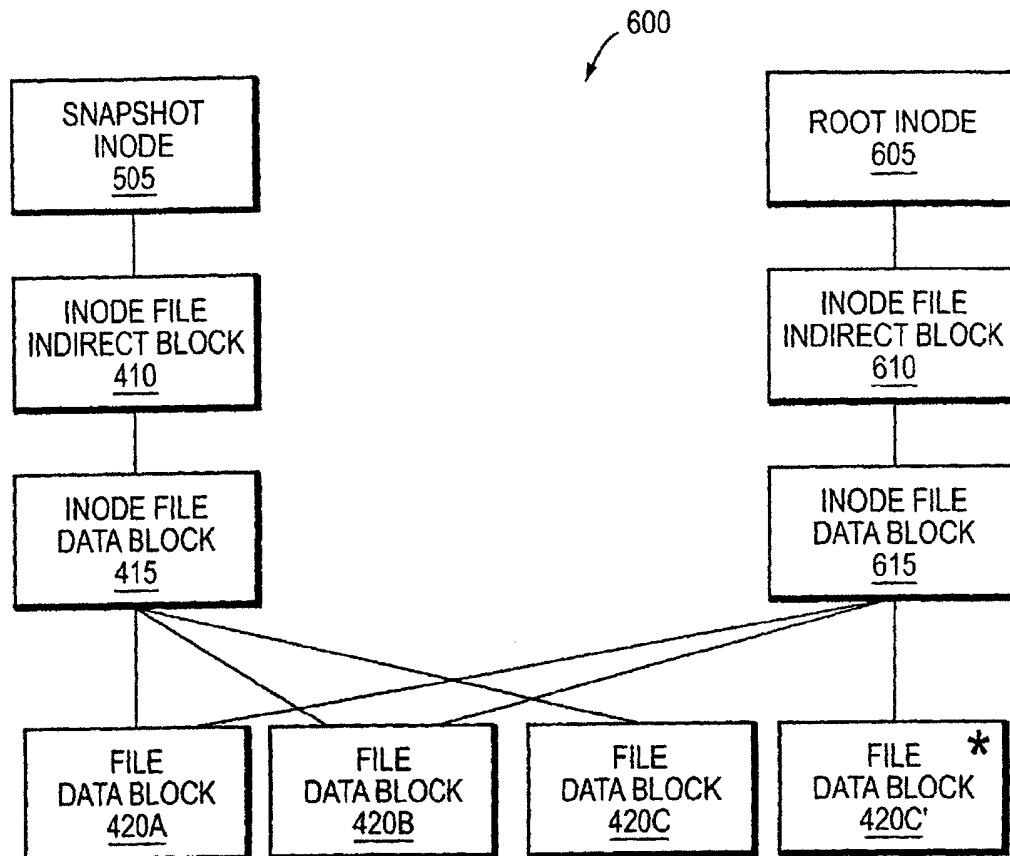


FIG. 6

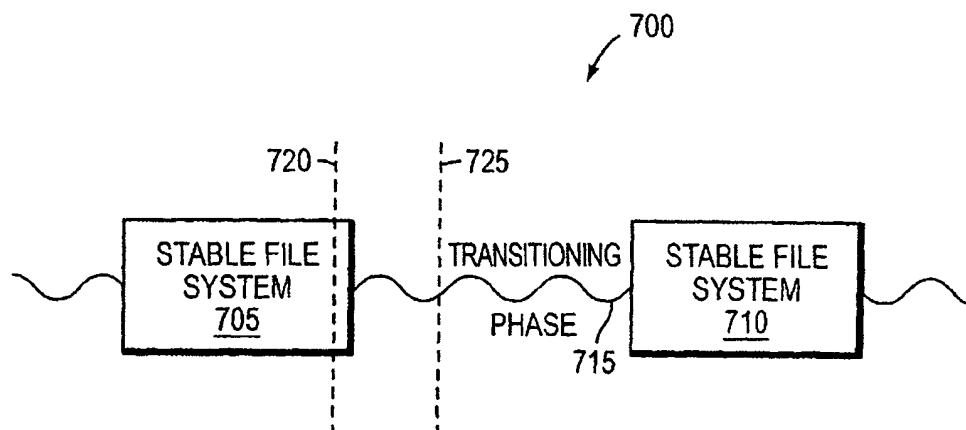


FIG. 7

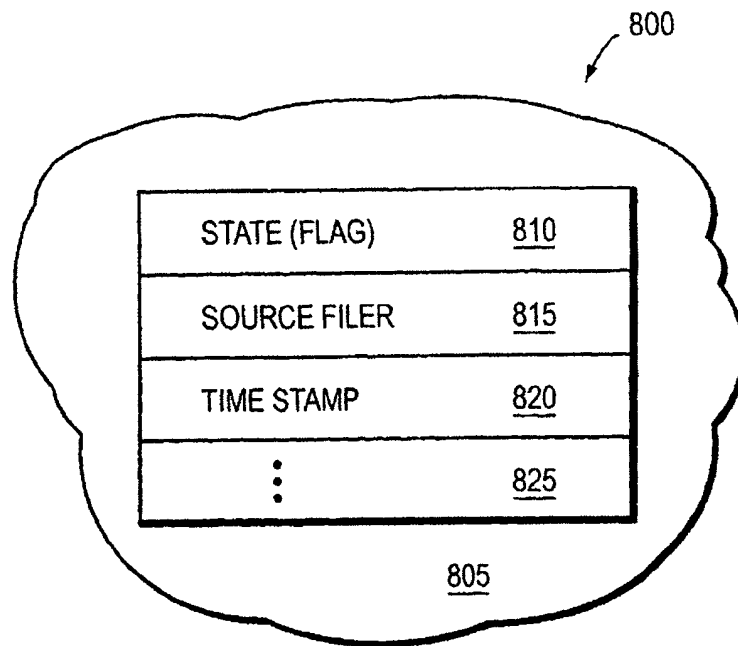


FIG. 8

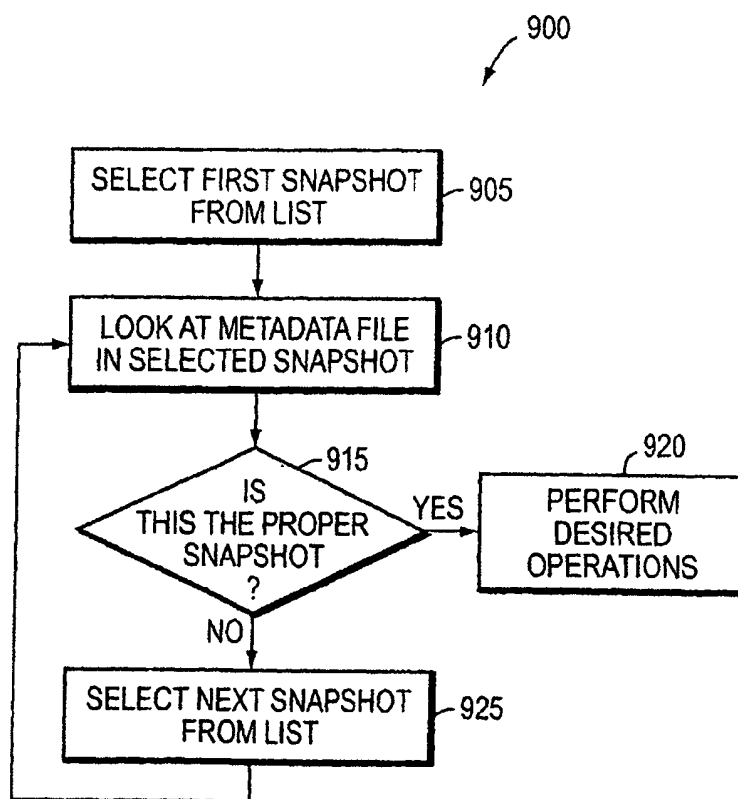


FIG. 9

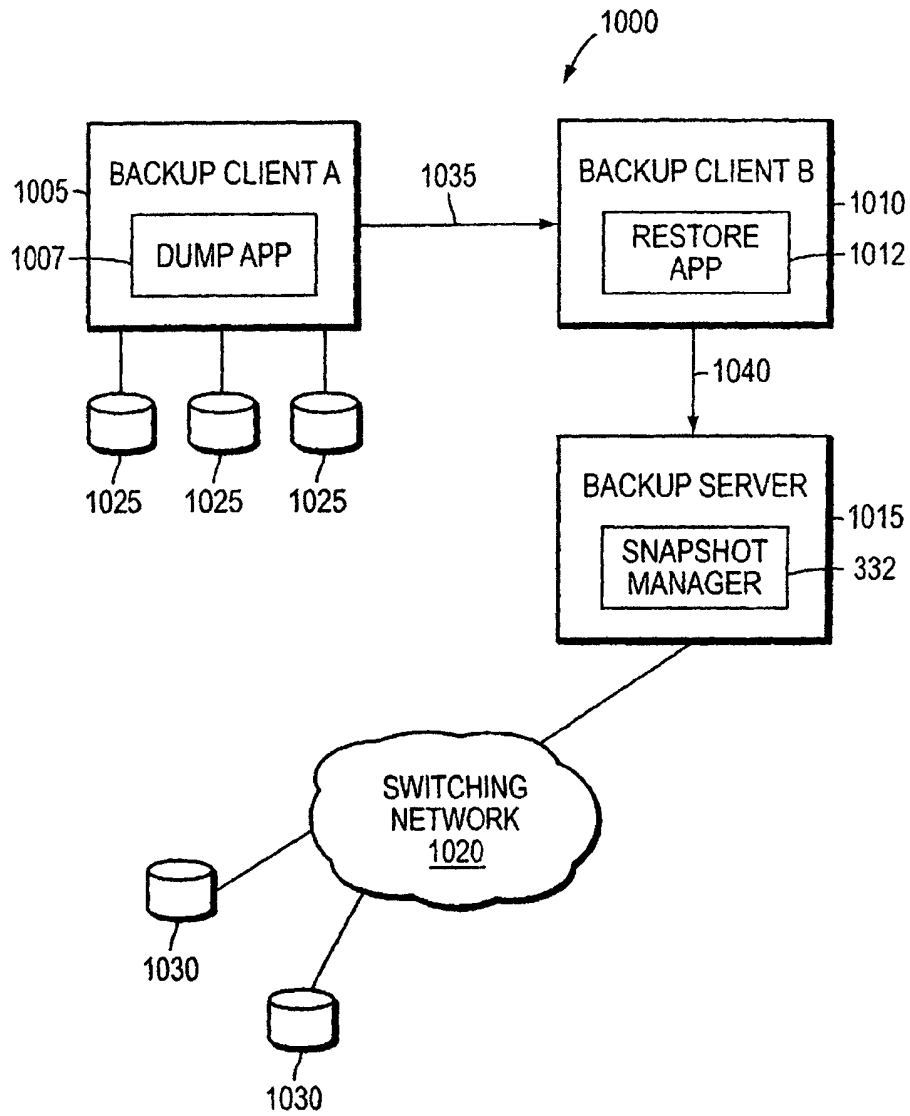


FIG. 10

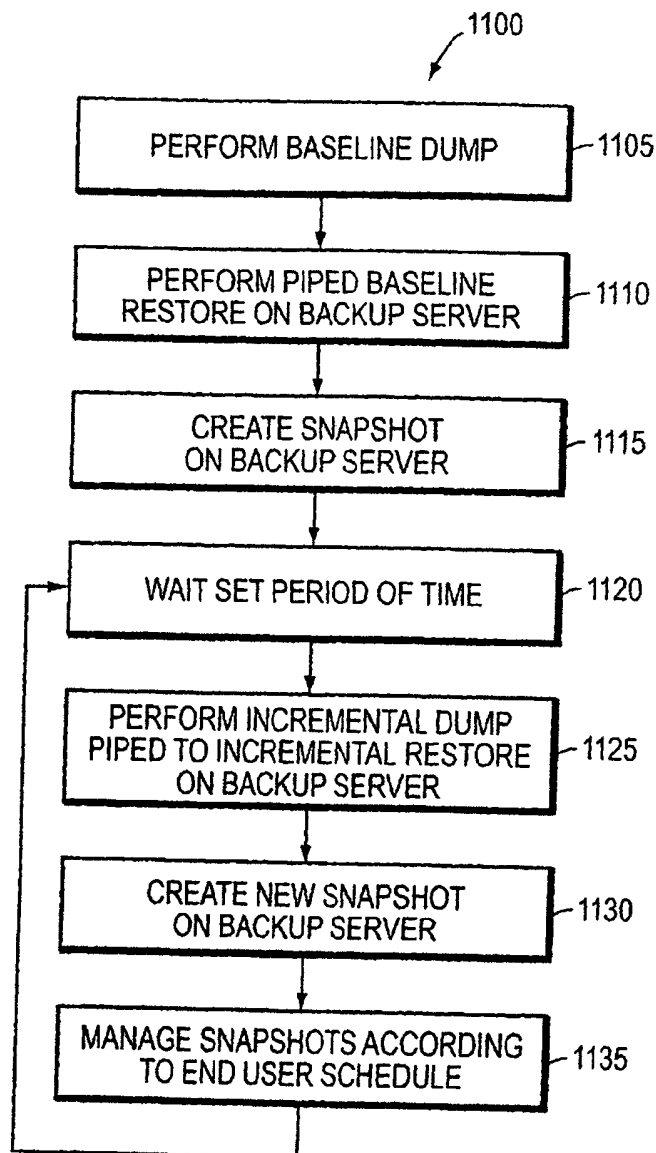


FIG. 11

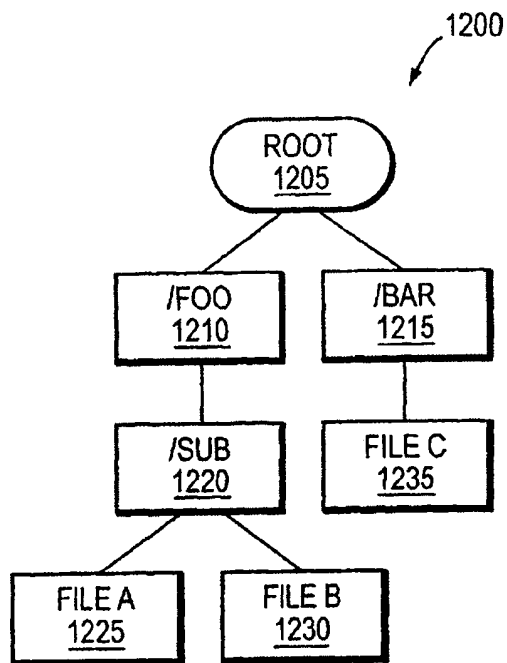


FIG. 12

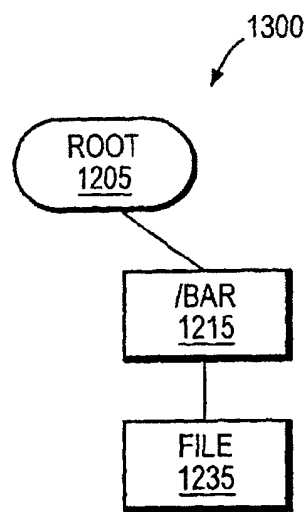
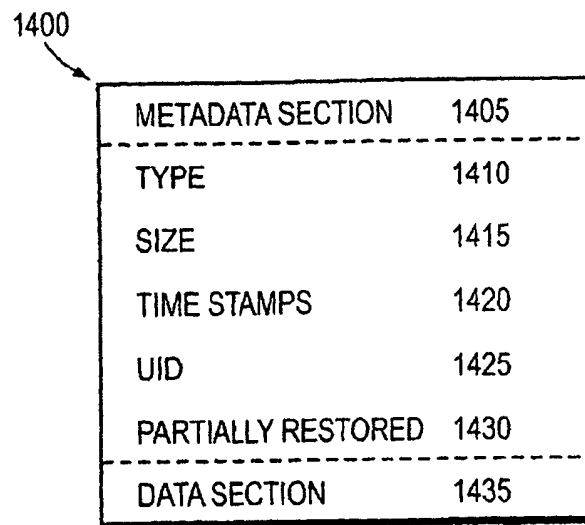


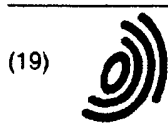
FIG. 13

1400



METADATA SECTION	1405
<hr/>	
TYPE	1410
SIZE	1415
TIME STAMPS	1420
UID	1425
PARTIALLY RESTORED	1430
<hr/>	
DATA SECTION	1435

FIG. 14



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(54) **System and method for managing a plurality of snapshots of a file system**

(57) A system and method for managing a plurality of snapshots as provided. A set of metadata describing a file system is contained within the file system so that a snapshot of the file system includes the associated metadata. Backup client file systems are restored to a backup server using conventional dump and restore techniques. The backup server then utilizes a user-defined snapshot management schedule to manage the set of backups associated with the backup server. Such management of snapshots can include deletion of snapshots based upon a variety of parameters including the timestamp.

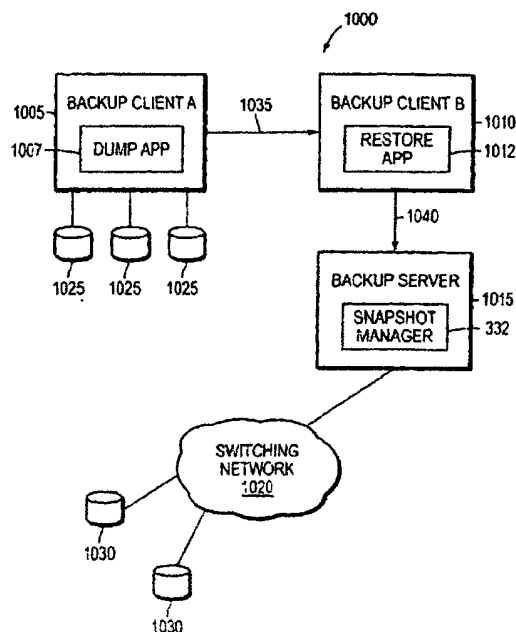


FIG. 10

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EUROPEAN SEARCH REPORT

Application Number
EP 03 25 1703

DOCUMENTS CONSIDERED TO BE RELEVANT			
Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	CLASSIFICATION OF THE APPLICATION (Int.Cl.7)
Y	US 5 835 953 A (OHRAN RICHARD) 10 November 1998 (1998-11-10) * abstract * * column 10, line 20 - column 11, line 47 * * column 12, line 35 - line 51 * * figure 3 *	1-12,20	G06F17/30 G06F11/14
Y	WO 00/07104 A (NETWORK APPLIANCE INC) 10 February 2000 (2000-02-10) * abstract * * page 8, line 14 - page 9, line 15 * * page 24, line 25 - page 25, line 26 * * figure 2 *	1-12,20	
A	WO 01/31446 A (NETWORK APPLIANCE INC) 3 May 2001 (2001-05-03) * abstract *	1-12,20	
P,A	US 2002/178146 A1 (AKELLA RAJI LAKSHMI ET AL) 28 November 2002 (2002-11-28) * abstract; figure 2 * * paragraphs [0006], [0007], [0025], [0037] *	1-12,20	TECHNICAL FIELDS SEARCHED (Int.Cl.7) G06F
<p>The present search report has been drawn up for all claims</p>			
Place of search Munich		Date of completion of the search 16 December 2004	Examiner Leuridan, K
<p>CATEGORY OF CITED DOCUMENTS</p> <p>X : particularly relevant if taken alone Y : particularly relevant if combined with another document of the same category A : technological background O : non-written disclosure P : intermediate document</p> <p>T : theory or principle underlying the invention E : earlier patent document, but published on, or after the filing date D : document cited in the application L : document cited for other reasons & : member of the same patent family, corresponding document</p>			

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EPO FORM 1303 03.02.02 (P04/C01)



European Patent
Office

Application Number
EP 03 25 1703

CLAIMS INCURRING FEES

The present European patent application comprised at the time of filing more than ten claims.

- ☐ Only part of the claims have been paid within the prescribed time limit. The present European search report has been drawn up for the first ten claims and for those claims for which claims fees have been paid, namely claim(s):
- ☐ No claims fees have been paid within the prescribed time limit. The present European search report has been drawn up for the first ten claims.

LACK OF UNITY OF INVENTION

The Search Division considers that the present European patent application does not comply with the requirements of unity of invention and relates to several inventions or groups of inventions, namely:

see sheet 8

- ☐ All further search fees have been paid within the fixed time limit. The present European search report has been drawn up for all claims.
- ☐ As all searchable claims could be searched without effort justifying an additional fee, the Search Division did not invite payment of any additional fee.
- ☐ Only part of the further search fees have been paid within the fixed time limit. The present European search report has been drawn up for those parts of the European patent application which relate to the inventions in respect of which search fees have been paid, namely claims:
- ☒ None of the further search fees have been paid within the fixed time limit. The present European search report has been drawn up for those parts of the European patent application which relate to the invention first mentioned in the claims, namely claims:

1-12, 20



European Patent
Office

**LACK OF UNITY OF INVENTION
SHEET B**

Application Number
EP 03 25 1703

The Search Division considers that the present European patent application does not comply with the requirements of unity of invention and relates to several inventions or groups of inventions, namely:

1. claims: 1-12,20

A method for efficient management of backups

2. claims: 13-17

A method of generating a self-describing snapshot

3. claims: 18,19

A method of accessing a file stored in a snapshot

**ANNEX TO THE EUROPEAN SEARCH REPORT
ON EUROPEAN PATENT APPLICATION NO.**

EP 03 25 1703

This annex lists the patent family members relating to the patent documents cited in the above-mentioned European search report.
The members are as contained in the European Patent Office EDP file on
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16-12-2004

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EPO FORM P0438

For more details about this annex : see Official Journal of the European Patent Office, No. 12/82